

Sensitivity analysis and polynomial chaos for differential equations

Dongbin Xiu

Department of Mathematics, Purdue University

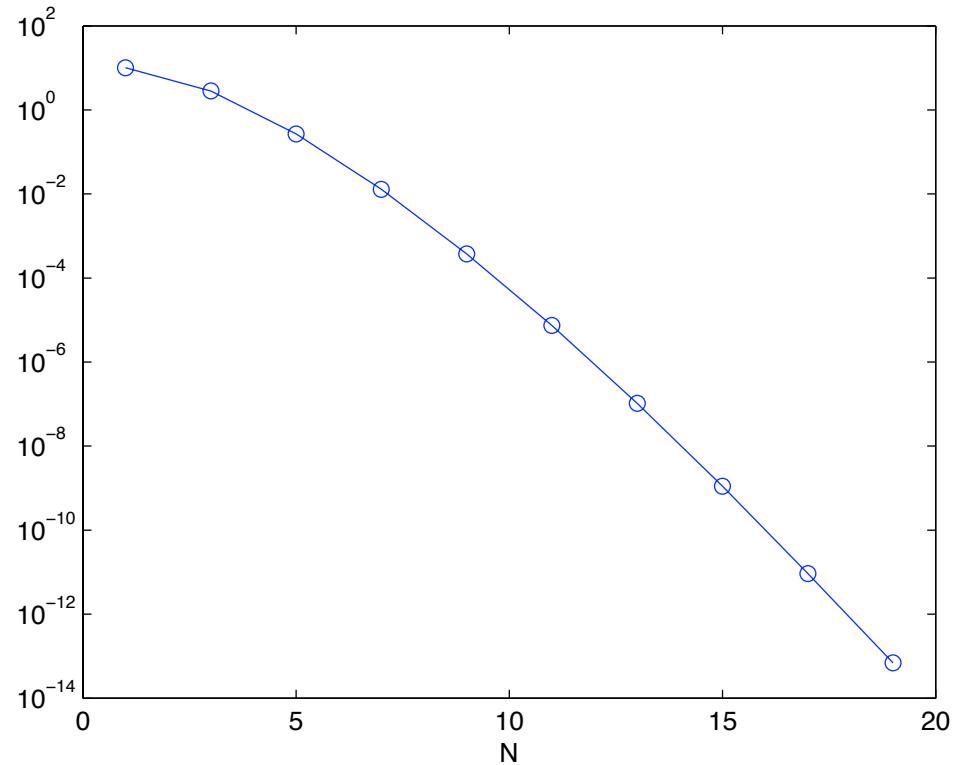


Support: AFOSR, DOE, NNSA, NSF

Outline

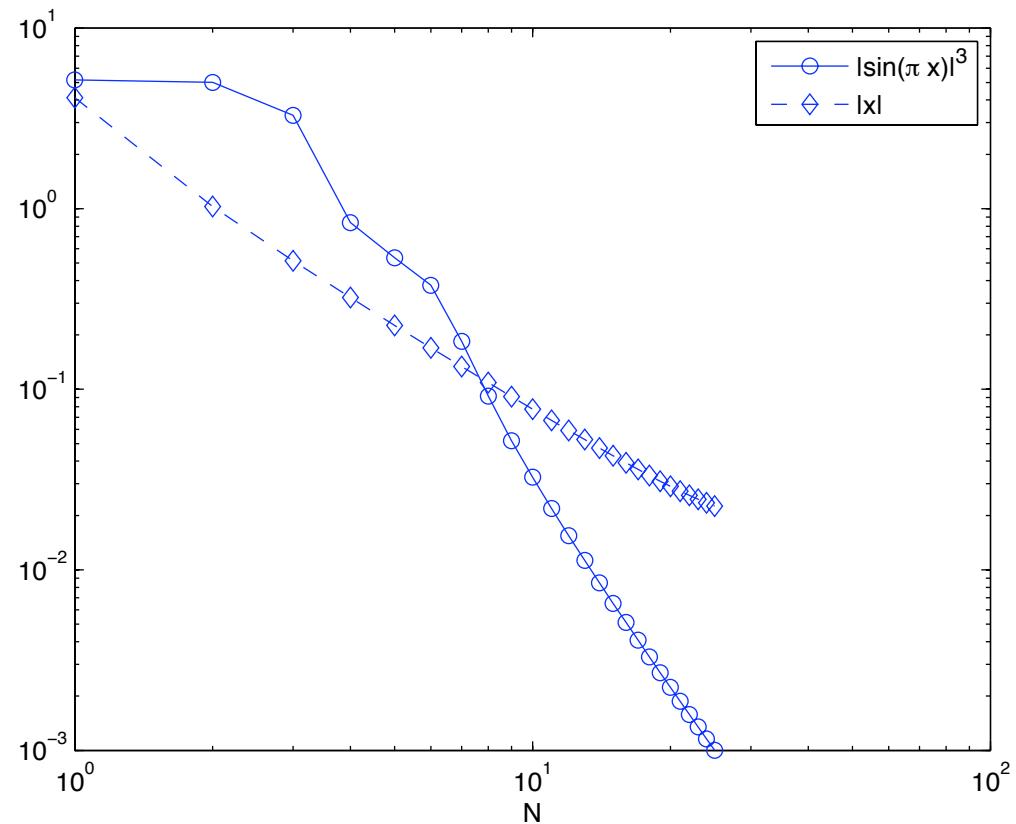
- Basic approximation theory
 - Formulation of SPDE
 - “Basic-er” probability facts
 - Sensitivity analysis
 - Generalized polynomial chaos
 - Numerical methods
 - Stochastic Galerkin
 - Stochastic collocation
- *Xiu, Numerical methods for stochastic computations, Princeton Univ. Press, 2010*

Spectral Accuracy: exponential convergence



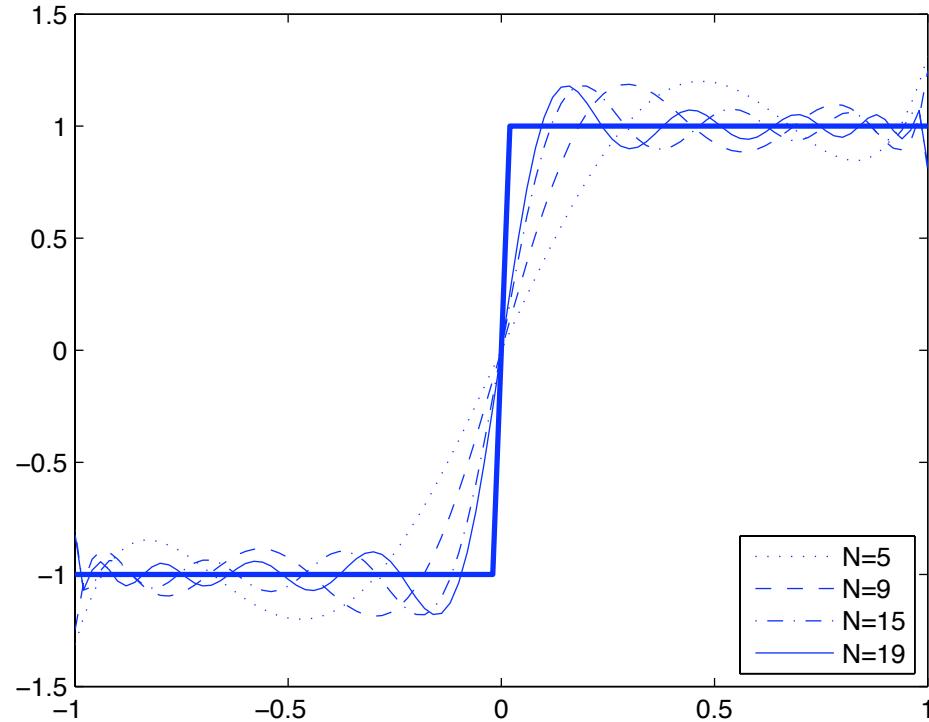
Projection error of $\cos(\pi x)$ by Legendre polynomials

Spectral Accuracy: algebraic convergence



Projection error of $|\sin(\pi x)|^3$ and $|x|$ by Legendre polynomials

Spectral Accuracy: Gibb's phenomenon



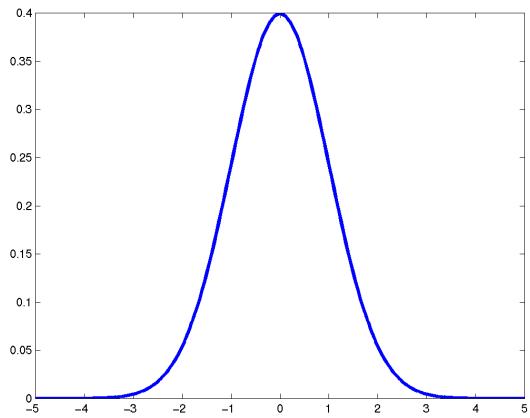
Legendre series expansion of sign function

$$\operatorname{sgn}(x) = \sum_{n=0}^{\infty} \frac{(-1)^n (4n+3)(2n)!}{2^{2n+1} (n+1)! n!} P_{2n+1}(x).$$

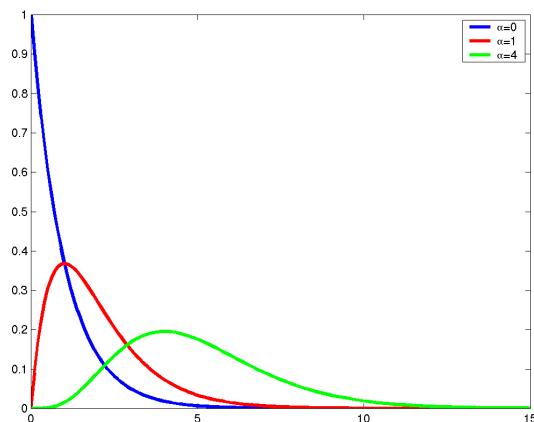
Probability distribution:

$$F_X(x) = P(X < x) \in [0,1]$$

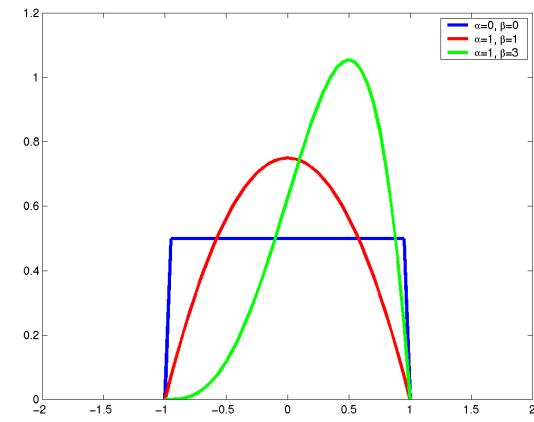
Probability density function (PDF): $\rho_X(x) = \frac{dF(x)}{dx}$



Gaussian distribution



Gamma distribution



Beta distribution

Expectation:

$$\mathbb{E}(g(X)) = \int g(x) dF_X(x) = \int g(x) \rho(x) dx$$

Mean: $\mu_X = \mathbb{E}(X)$

Variance: $\sigma^2 = \mathbb{E}((X - \mu_X)^2)$

Skewness: $\mathbb{E}((X - \mu_X)^3)$

Kurtosis: $\mathbb{E}((X - \mu_X)^4)$

(Re-)Formulation of PDE: Input Parameterization

$$\frac{\partial u}{\partial t}(t, x) = \mathcal{L}(u) + \text{boundary/initial conditions}$$

- **Goal:** To characterize the random inputs by a set of random variables
 - Finite number
 - Mutual independence
- **If inputs == parameters**
 - Identify the (smallest) independent set
 - Prescribe probability distribution
- **Else if inputs == fields/processes**
 - Approximate the field by a function of finite number of RVs
 - Well-studied for Gaussian processes
 - Under-developed for non-Gaussian processes
 - Examples: Karhunen-Loeve expansion, spectral decomposition, etc.

$$a(x, \omega) \approx \mu_a(x) + \sum_{i=1}^d \tilde{a}_i(x) Z_i(\omega)$$

The Reformulation

- **Stochastic PDE:**

$$\frac{\partial u}{\partial t}(t, x, Z) = \mathcal{L}(u) + \text{boundary/initial conditions}$$

- **Solution:** $u(t, x, Z) : [0, T] \times \bar{D} \times \mathbb{R}^{n_z} \rightarrow \mathbb{R}$

- Uncertain inputs are characterized by n_z random variables Z

- Probability distribution of Z is prescribed

$$F_Z(s) = \Pr(Z \leq s), \quad s \in \mathbb{R}^{n_z}$$

Non-trivial task

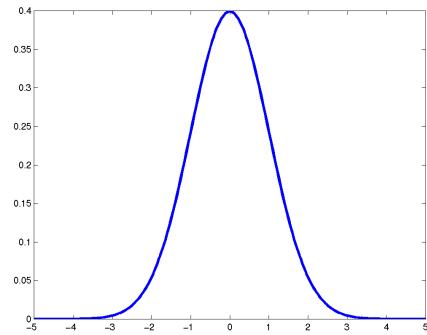
gPC Basis

- **Expectation:**

$$\mathbb{E}(g(Z)) = \int_{\mathbb{R}} g(z)\rho(z)dz$$

- **Orthogonality:**

$$\int \Phi_i(z)\Phi_j(z)\rho(z)dz = \mathbb{E}[\Phi_i(Z)\Phi_j(Z)] = \delta_{ij}$$

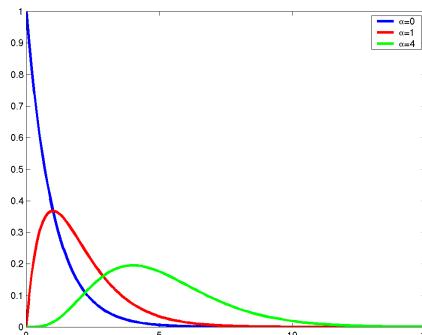


Gaussian distribution

$$\int_{-\infty}^{\infty} \Phi_i(z)\Phi_j(z)e^{-z^2} dz = \delta_{ij}$$



Hermite polynomial

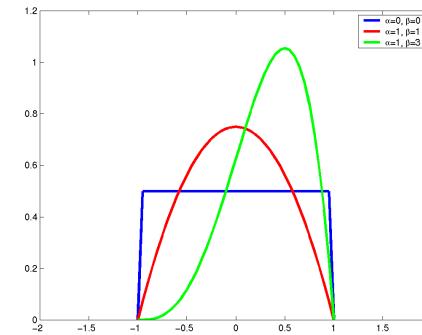


Gamma distribution

$$\int_0^{\infty} \Phi_i(z)\Phi_j(z)e^{-z} dz = \delta_{ij}$$



Laguerre polynomial



Beta distribution

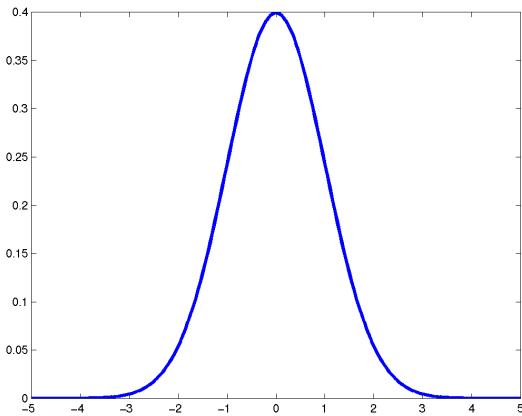
$$\int_{-1}^1 \Phi_i(z)\Phi_j(z)dz = \delta_{ij}$$



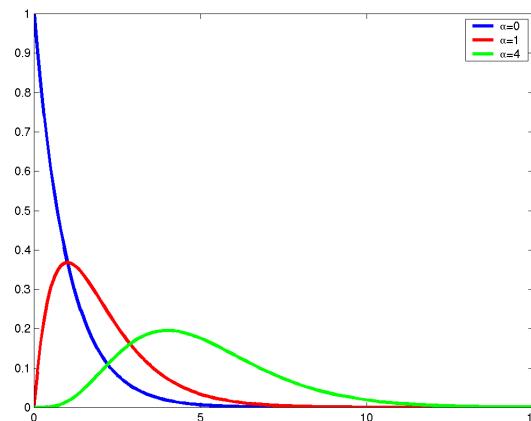
Legendre polynomial

■ Continuous Cases:

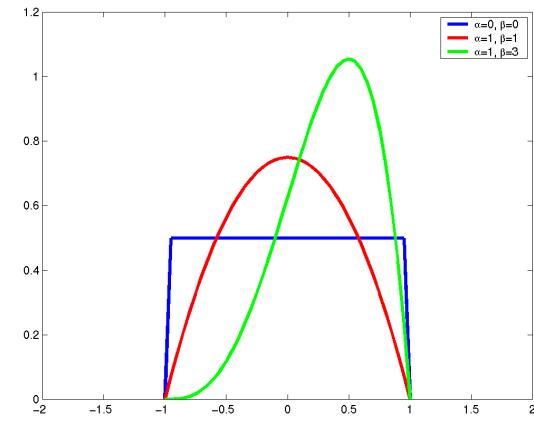
- *Hermite Polynomials* \longleftrightarrow *Gaussian Distribution*
- *Laguerre Polynomials* \longleftrightarrow *Gamma Distribution*
(special case: *exponential distribution*)
- *Jacobi Polynomials* \longleftrightarrow *Beta Distribution*
- *Legendre Polynomials* \longleftrightarrow *Uniform Distribution*



Gaussian distribution



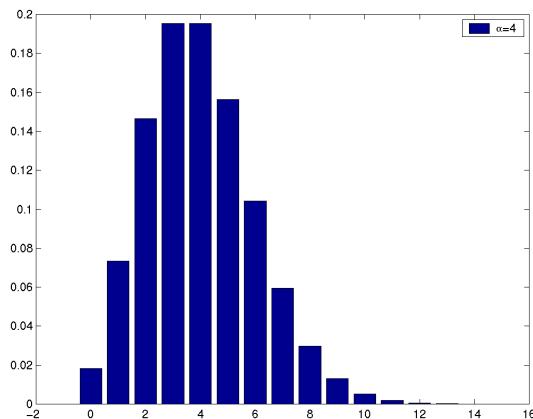
Gamma distribution



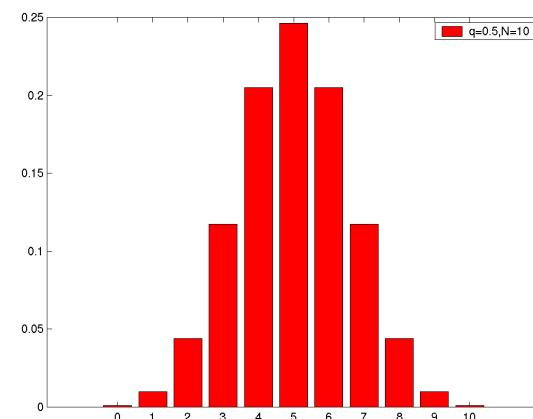
Beta distribution

■ Discrete Cases :

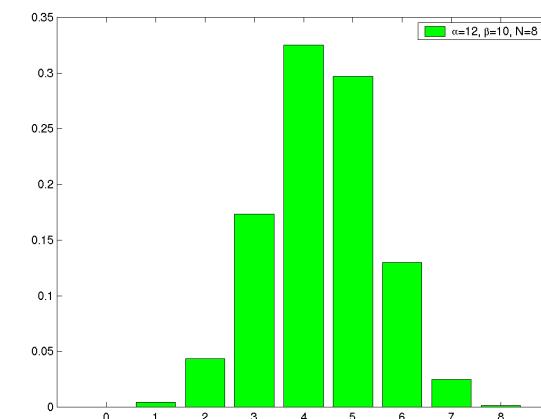
- *Charlier Polynomials* \longleftrightarrow *Poisson Distribution*
- *Krawtchouk Polynomials* \longleftrightarrow *Binomial Distribution*
- *Hahn Polynomials* \longleftrightarrow *Hypergeometric Distribution*
- *Meixner Polynomials* \longleftrightarrow *Pascal Distribution*



Poisson distribution



Binomial distribution

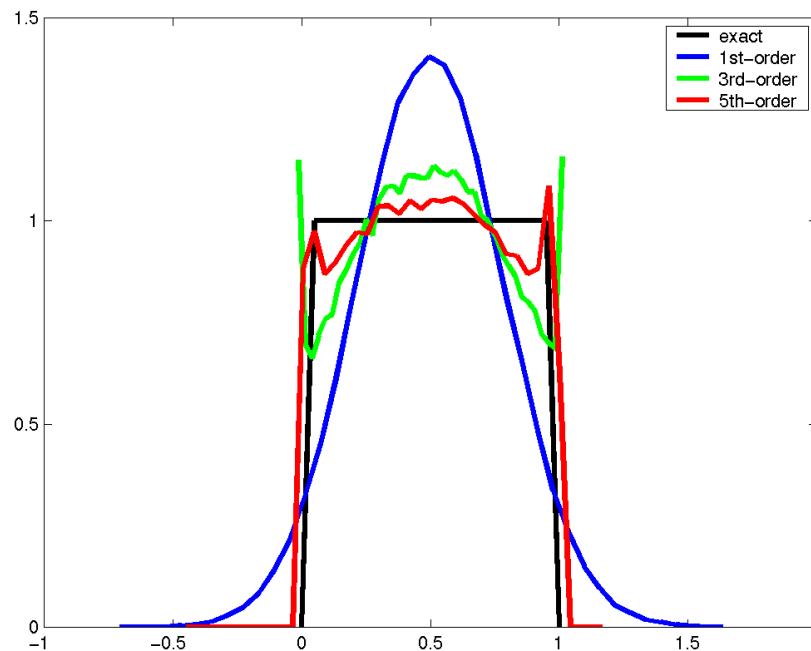


Hypergeometric distribution

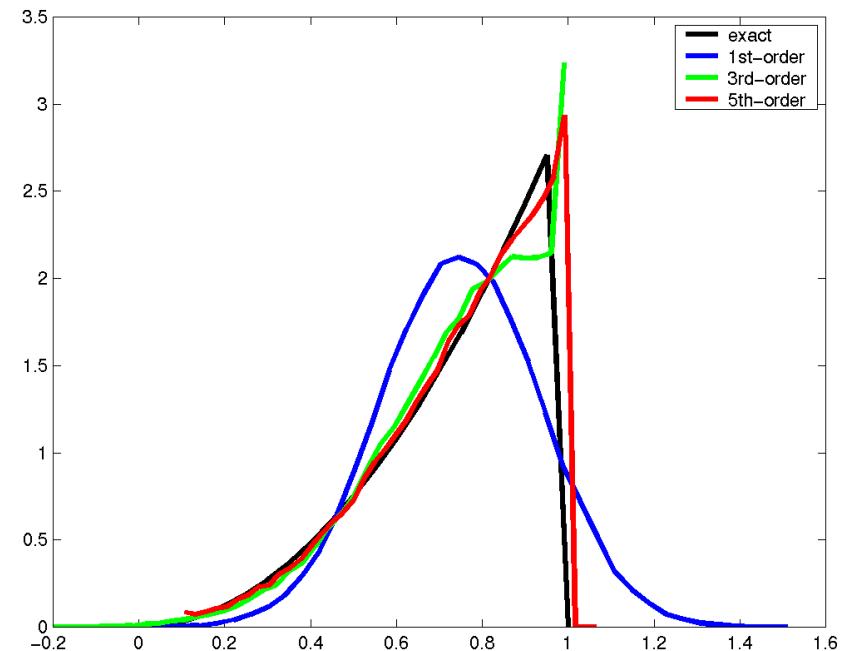
Hermite-Chaos Expansion of Beta Distribution

$$\text{PDF: } f(x) = \frac{x^{\alpha-1}(1-x)^{\beta-1}}{B(\alpha, \beta)}, \quad \alpha, \beta > 0, \quad 0 \leq x \leq 1$$

Uniform distribution : $\alpha = 1, \beta = 1$



$\alpha = 3, \beta = 1$



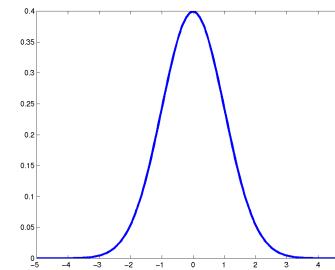
Exact PDF and PDF of 1st, 3rd, 5th-order Hermite-Chaos Expansions

gPC Basis: the Choices

- Orthogonality: $\int \Phi_i(z)\Phi_j(z)\rho(z)dz = \mathbb{E}[\Phi_i(Z)\Phi_j(Z)] = \delta_{ij}$

- Example: Hermite polynomial

$$\int_{-\infty}^{\infty} \Phi_i(z)\Phi_j(z)e^{-z^2} dz = \delta_{ij}$$



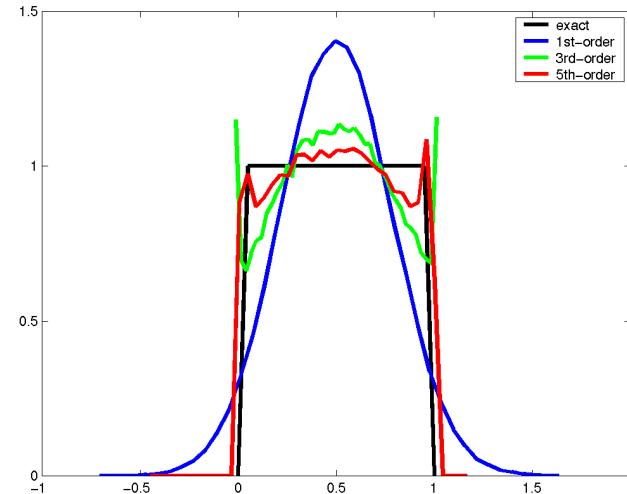
- The polynomials: $Z \sim N(0, 1)$

$$\Phi_0 = 1, \quad \Phi_1 = Z, \quad \Phi_2 = Z^2 - 1, \quad \Phi_3 = Z^3 - 3Z, \quad \dots$$

- Approximation of arbitrary random variable: Requires L^2 integrability

- Example: Uniform random variable

- Convergence
- Non-optimal
- First-order Legendre is exact



Stochastic Galerkin

$$\frac{\partial u}{\partial t}(t, x, Z) = \mathcal{L}(u) + \text{boundary/initial conditions}$$

- **Galerkin method:** Seek

$$u_N(t, x, Z) \triangleq \sum_{|\mathbf{k}|=0}^N \hat{u}_{\mathbf{k}}(t, x) \Phi_{\mathbf{k}}(Z)$$

Such that

$$\mathbb{E}\left[\frac{\partial u_N}{\partial t}(t, x, Z) \Phi_{\mathbf{m}}(Z)\right] = \mathbb{E}\left[\mathcal{L}(u_N) \Phi_{\mathbf{m}}(Z)\right], \quad \forall |\mathbf{m}| \leq N$$

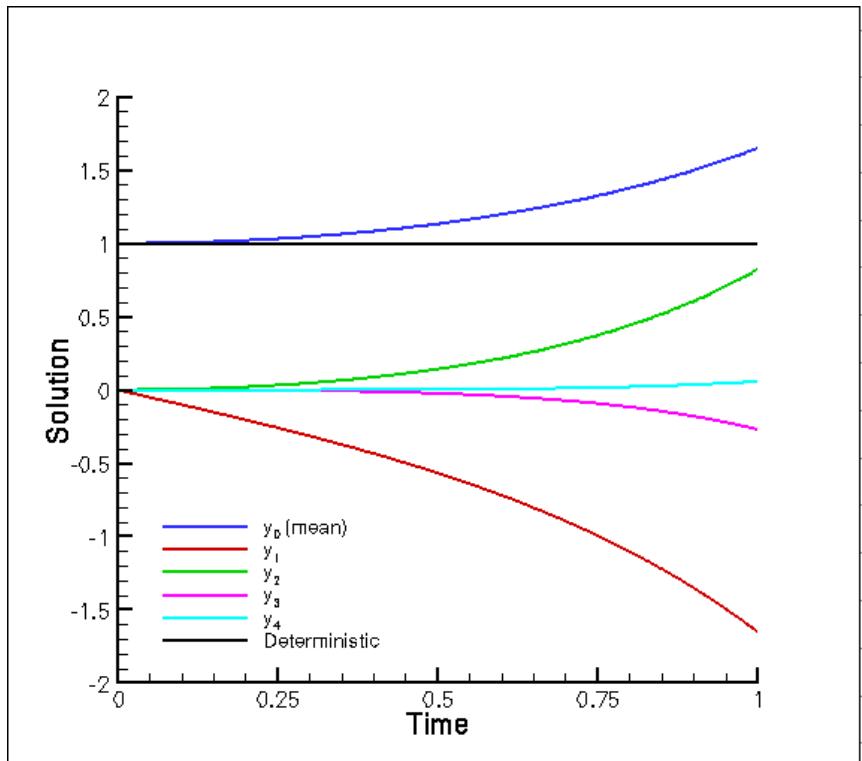
- **The result:**

- Residue is orthogonal to the gPC space
- A set of deterministic equations for the coefficients
- The equations are usually coupled – requires new solver

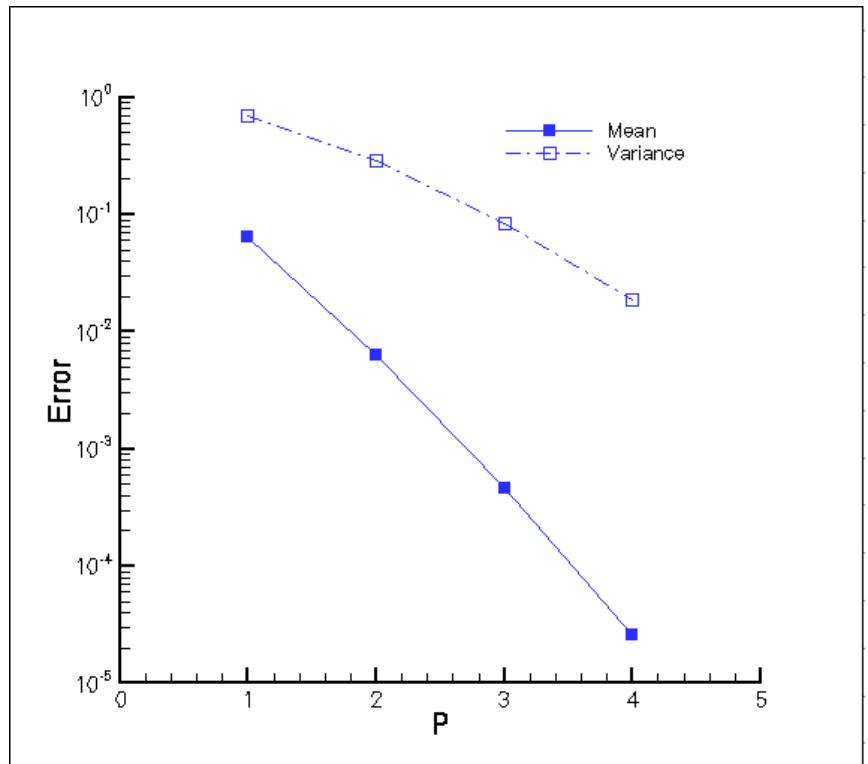
Continuous Distribution : Gaussian (Hermite-Chaos)

- $du/dt = - Z y, \ u(t=0)=1$
- Z is a **Gaussian** random variable :

$$\text{PDF: } f_Z(z) = \frac{1}{\sqrt{2\pi}} e^{-\frac{z^2}{2}}$$



Solution of expansion coefficients



Convergence w.r.t. expansion terms

- 4th-order **Hermite-Chaos** expansion
- Exponential convergence rate

Stochastic Collocation

$$\frac{\partial u}{\partial t}(t, x, Z) = \mathcal{L}(u) \quad + \text{boundary/initial conditions}$$

- **Collocation:** To satisfy governing equations at selected nodes
 - Allow one to use existing deterministic codes repetitively
-

- **Sampling:** (solution statistics only)
 - Random (Monte Carlo)
 - Deterministic (lattice rule, tensor grid, cubature)
-

- **Stochastic collocation:** To construct **polynomial approximations**
 - Node selection is critical to efficiency and accuracy
 - More than sampling

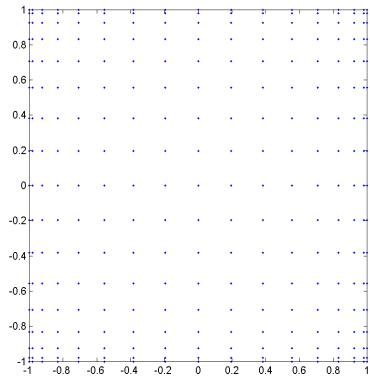
Stochastic Collocation: Interpolation

$$\frac{\partial u}{\partial t}(t, x, Z) = \mathcal{L}(u) + \text{boundary/initial conditions}$$

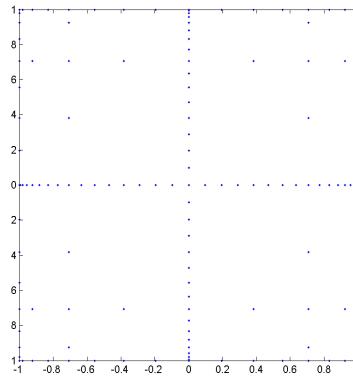
- **Definition:** Given a set of nodes and solution ensemble, find u_N in a proper polynomial space, such that $u_N \approx u$ in a proper sense.

- **Interpolation Approaches:** $u_N(Z) = \sum_{j=1}^Q u(Z^j) L_j(Z)$
 $L_i(Z^j) = \delta_{ij}, \quad 1 \leq i, j \leq Q$

- Dimension-by-dimension space filling



Tensor grids: inefficient



Sparse grids: more efficient

gPC-Collocation: Algorithm

Deterministic solver

1. Choose a nodal set $\{Z^j, \alpha^j\}_{j=1}^Q$ in \mathbb{R}^{n_z}

2. Solve for each $j = 1, \dots, Q$,

$$\frac{\partial u}{\partial t}(t, x, Z^j) = \mathcal{L}(u), \quad \text{in } (0, T] \times D,$$

$$\mathcal{B}(u) = 0, \quad [0, T] \times \partial D,$$

$$u = u_0(x, Z^j), \quad \{t = 0\} \times D$$

3. Evaluate the approximate gPC expansion coefficient

$$\hat{w}_k = \sum_{j=1}^Q u(t, x, Z^j) \Phi_k(Z^j) \alpha^j, \quad 0 \leq |k| \leq N;$$

Post-process

4. Construct the N^{th} -order gPC approximation

$$w_N(t, x, Z) = \sum_{|k|=1}^N \hat{w}_k \Phi_k(Z).$$

• **Error bound:**

$$\varepsilon \triangleq \|u - w_N\|_{L^2(Z)} \leq (\varepsilon_N^2 + \varepsilon_Q^2 + M \varepsilon_\Delta^2 C_Q^2)^{1/2}$$

Error \leq

Finite-term projection error

aliasing error

Numerical error